## Architecture and Performance characteristics of a PostgreSQL implementation of the TPC-E and TPC-V workloads

Andrew Bond (Red Hat), Doug Johnson (InfoSizing), Greg Kopczynski (VMware), and H. Reza Taheri (VMware)



## **Agenda/Topics**

- Virtualization and virtualization benchmarks
- Historical perspective of TPC-V
- TPC-V design considerations and characteristics
- TPC-V architecture
- End-to-end Reference Kit
- Prototyping results, and PostgreSQL characterization
- Benchmark roadmap

### Virtualization and need for a database benchmark

#### Virtualizing servers allows:

- consolidation
  - Reduces both CapEx and OpEx
- Migration of VMs (both storage and CPU/memory)
  - Enables live load balancing
  - Facilitates maintenance
- High availability and fault tolerance
- Cloud computing is powered by virtualized servers
- Databases VMs are were the last frontier for virtualization
  - Initial hesitation to put enterprise databases on VMs
  - Followed by today's push towards virtualizing everything
- We need a benchmark to model all this!

### **Today's virtualization benchmarks**

### VMmark

- Developed by VMware in 2007; now on version 2.5
- De facto industry standard with nearly 200+ publications from 11 vendors
- Released kit only runs on ESX, but can be modified for other hypervisors
- Models *consolidation* of lightly loaded VMs with diverse workloads

### SPECvirt\_sc2010/SPECvirt\_sc2013

- Industry standard, with 35 results from 4 vendors
- Models *consolidation* of lightly loaded VMs with diverse workloads

### TPC-VMS

- No results yet
- Models consolidation of 3 identical database VMs
  - Test sponsors can use one of 4 existing TPC benchmark workloads
  - Limited in scope of what virtualization features get tested
  - +But does not require a *new* kit

### **History of TPC-V**

#### 2009 VLDB: We need a benchmark that:

- Models virtualized databases
- Is industry standard

#### 2010 VLDB: Proposal for a benchmark

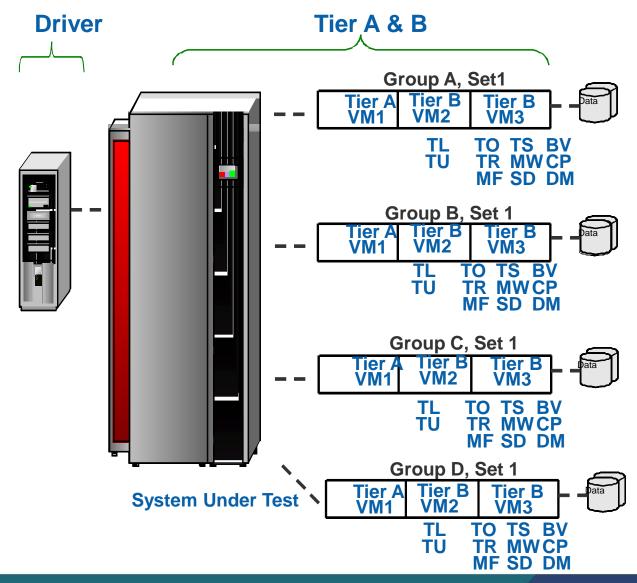
- Steal TPC-E's workload, schema, specification
- Baby bear # of VMs
- Model elasticity of load

### 2012 VLDB: Status of the benchmark development

- Developing a complete, end-to-end, publicly-available kit
- Kit runs on PostgreSQL
- Finalized benchmark architecture
- Some early result

This pace is lightening fast by TPC standards!

### **Components of a TPC-V configuration**

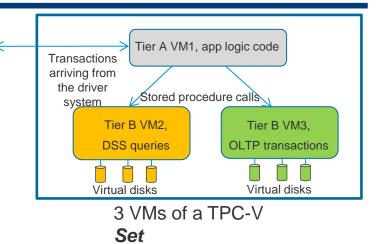


**VLDB 2013** 

### **Sets and Groups**

Group:	Α	В	С	D
Average contribution	10%	20%	30%	40%
to overall throughout				

- VMs with heterogeneous load levels
  - Always 4 *Groups* @10%, 20%, 30%, 40%
- But the number of Sets per Group grows with performance
  - Small systems have 1 Set per Group
  - Larger servers divide each Group's load among 2, 3, ... Sets
- Total number of VMs on the server:
  - 12 on small and medium-sized servers
  - 24 on today's high end
  - 36-48 on high end in 5-10 years



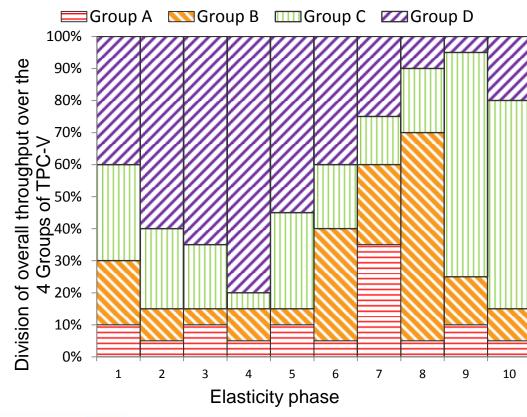
Group A, Set 1	VM1 A1	VM2 A1	VM3 A1
Group B, Set 1	VM1 B1	VM2 B1	VM3 B1
Group C, Set 1	VM1 C1	VM2 C1	VM3 C1
Group D, Set 1	VM1 D1	VM2 D1	VM3 D1

VLDB 2013

### Elasticity



- Elasticity
- Oversubscription

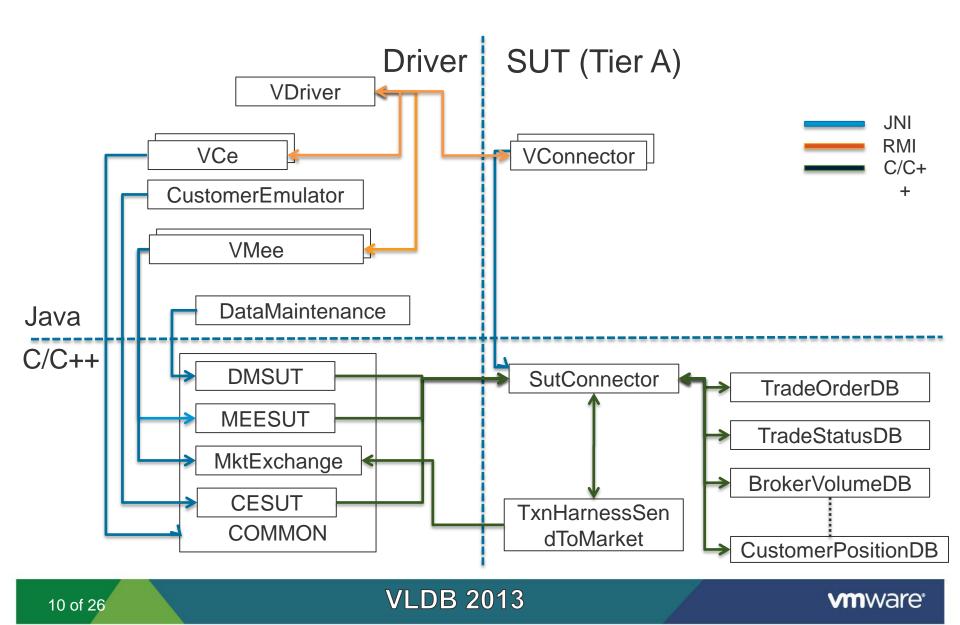


Period	Group A	Group B	Group B	Group D
1	10%	20%	30%	40%
2	5%	10%	25%	60%
3	10%	5%	20%	65%
4	5%	10%	5%	80%
5	10%	5%	30%	55%
6	5%	35%	20%	40%
7	35%	25%	15%	25%
8	5%	65%	20%	10%
9	10%	15%	70%	5%
10	5%	10%	65%	20%
Average	10%	20%	30%	40%

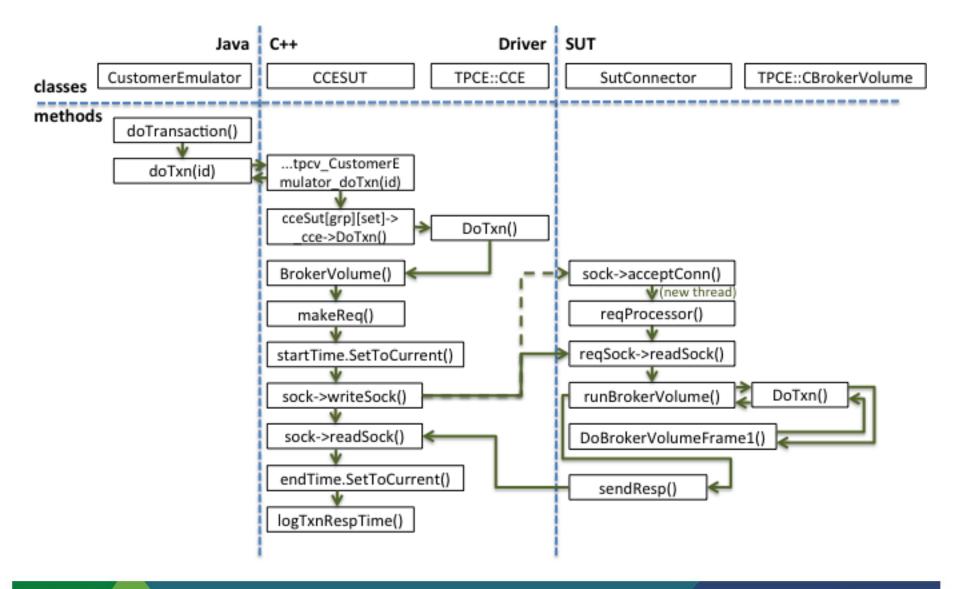
**VLDB 2013** 

### **Design Considerations**

- Driver code in Java; transaction-specific code in C++ to match EGen/VGen
- Walk before you run: a TPC-E kit first
- ODBC allows easier database swapping
  - Albeit with a performance cost
  - We can replace ODBC with native calls for best performance
- Develop initial kit on PostgreSQL
  - But ODBC makes it possible to run against other databases
- Make the kit available to anyone, subject to a EULA



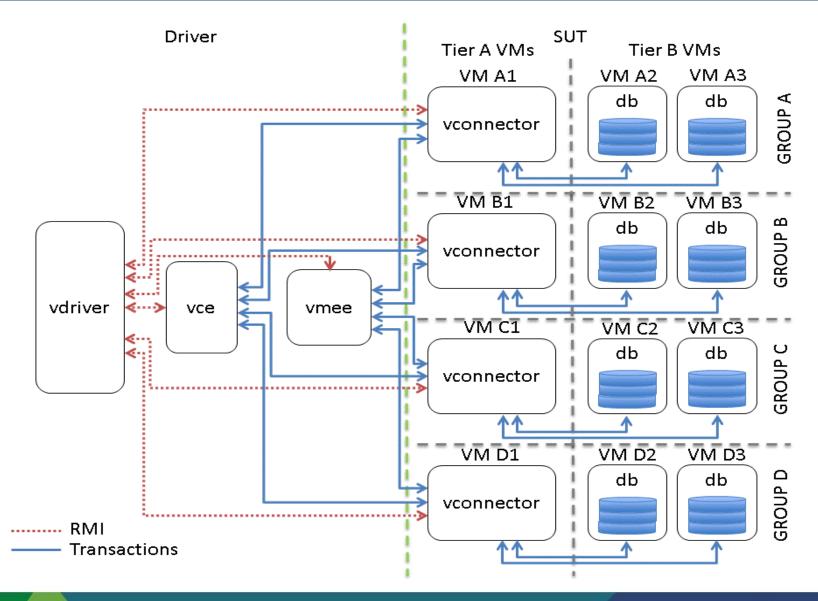
### **Request Code Execution Path**



**VLDB 2013** 

### **Driver components**

12 of 26



**VLDB 2013** 

### Status of the kit

#### We have a complete, end-to-end TPC-E kit

- This is TPC-E, so 1 database, static load
- Scripts to create the schema and populate the database
- All the necessary DDL/DML
- C++ code to implement the business logic
- Java driver to drive the load and collect and process the results
  - Handles multiple Sets and Groups, and elasticity
- Linux shell scripts to start and monitor the run, and collect stats
- The kit implements all the multi-VM and elasticity features of TPC-V
  - Need to change the code for 2 transactions that are different in TPC-V
- Busy prototyping!

### **TPC-E results with a single database VM**

- Single VM/database (emulating TPC-E)
- 4-socket HP ProLiant DL580 G7 server
  - 2.40GHz Intel Xeon E7-4870 (WestmereEX) CPUs
    - So total of 40 cores/80 hyperthreads
  - 1TB of memory
  - Disclosed TPC-E result: 2,454 tpsE
    - Other results on similar servers as high as 3,XXX tpsE
  - Two EMC VNX5700 disk arrays. 38 SSDs and 88 spinning disk drives

### Running TPCE on a VM with 16-vCPU, 280GB Tier B VM

• So using 1/5<sup>th</sup> of the hardware resources

#### Software versions:

- VMware vSphere 5.1
- RHEL 6.1
- PGSQL 9.2.2
- unixODBC 2.2.14

**VLDB 2013** 

### **TPC-E prototyping results**

#### Measuring 198 tpsE on the 16-vCPU VM

- Performance is decent for this stage of development
  - 1/13<sup>th</sup> of disclosed results on this server, but using 1/5<sup>th</sup> the resources
    - So our performance is 1/3 to 1/2 of commercial databases
  - 2 orders of magnitude higher than dbt5 experimental results
- 85% CPU utilization
- 19K IOPS, 212MB/sec

#### I/O rate appears to be our biggest problem

- High IOPS even though we have cut the DB size by 60%
- IOPS/tran around 2X what commercial database does
- PostgreSQL buffer management and file system relationship unknowns
- Lack of clustered indexes
- On-disk size appears as much as 20% larger

### **PostgreSQL tuning: File systems**

- This data is from when we separated the OLTP and DSS transactions into two different VMs
- Definitely should separate the log and data file systems
  - So data blocks don't get flushed every time we write to the log
  - 6% more throughput
  - Lower response times

#### Use ext4 for data

• ext3 is fine for log

OLTP Trans		wrqm/s	r/s	w/s	rkB/s	wkB/s	avgrq- sz	avgqu- sz	await
1 file system	Data +log	403	542	476	7682	5552	27	5.1	4.75
2 file systems	Data	194	860	145	15613	1357	34	6.3	6.29
	log	1	0.04	225	0.16	3066	27	0.3	1.15
DSS Trans		wrqm/s	r/s	w/s	rkB/s	wkB/s	avgrq- sz	avgqu- sz	await
1 file system	Data +log	1830	11151	2767	138602	33956	25	30	2.14
2 file systems	data	2406	12350	2278	181902	18737	27	40	2.71
2 me systems	log	343	0.34	134	1	17854	264	0.3	1.87
16 of 26	VLDB 2013				VI	<b>n</b> ware <sup>®</sup>			

### **PostgreSQL** tuning: File systems, continued

#### PostgreSQL folks like to encourage you to use the file system

- No! They insist that relying on the filsys buffer cache is the only way to go
- Being an old database hand, you have to show me why this is so
  - Double buffering wastes memory
  - Once you miss in the DBMS cache, you pay most of the price in OS cycle whether or not you hit in the buffer cache

#### I have been trying to max out PostgreSQL shared\_buffers

- Ignoring the PostgreSQL book advice
- Hard to shrink file system buffer cache when raising shared\_buffers
  - End up swapping

17 of 26

- I have to reboot the VM
- But going from 34GB of shared\_buffers to 200GB improves performance only 5%
  - PostgreSQL/filsys interaction problem?

PC-E benefits from DBWSDBc2013 aching diminishing returns mware

### **PostgreSQL tuning: checkpoints**

- Watch the number of 16MB checkpoint\_segments
  - On a high throughput system, this will decide checkpoint frequency
  - **Default of only 3** checkpoint\_segments way too low
  - Even with 128 checkpoint\_segments, we were checkpointing every 2 minutes
  - Raised to 1,920 on the largest VM to checkpoints every 30 minutes

Checkpoint metric	128	segments	5,120 segments
checkpoints_timed		0	1
checkpoints_req		15	0
buffers_checkpoint		4,437,177	956,174
buffers_clean		14,069	852,893
buffers_backend		46,297	39,297
buffers_alloc		24,831,473	23,749,499

### **PostgreSQL tuning: Index-only scans**

#### TPC-E transactions need clustered indexes

- All published TPC-E results have used MS SQL Server
- They create clustered indexes on ALL 33 tables
- (Secondary) indexes not used very often
- PostgreSQL does not have clustered indexes
  - Every data access has to read both the index block and the data block
  - Indexes are not much smaller than tables
  - Plus, PostgreSQL on-disk footprint is larger

#### Switched to Index-only scans with PostgreSQL 9

- But works only if all the columns named in the query are in the index
- So have to create many, multi-column indexes
- Index size grows quickly
  - After a week of runs, TRADE data blocks grew 4%, index blocks 40%

### **PostgreSQL tuning: Index-only scans, continued**

Table		MS SQL		PostgreSQL			
	Table size	Index size	# of indexes	Table size	Index size	# of indexes	
CASH_TRANSACTION	125GB	0.45GB	1	146GB	140GB	1	
DAILY_MARKET	8.7GB	3.6GB	1	11GB	4.7GB	1	
HOLDING_HISTORY	67GB	35GB	2	93GB	125GB	2	
NEWS_ITEM	21GB	0.0003GB	1	20GB	0.003GB	1	
SETTLEMENT	68GB	0.3GB	1	91GB	78GB	1	
TRADE	153GB	82GB	3	176GB	135GB	3	
TRADE_HISTORY	96GB	0.25GB	1	168GB	124GB	1	

#### Note the index size for TRADE\_HISTORY

• Even having Index-only scans doesn't reduce how much data we juggle

### Server configuration for TPC-V prototyping

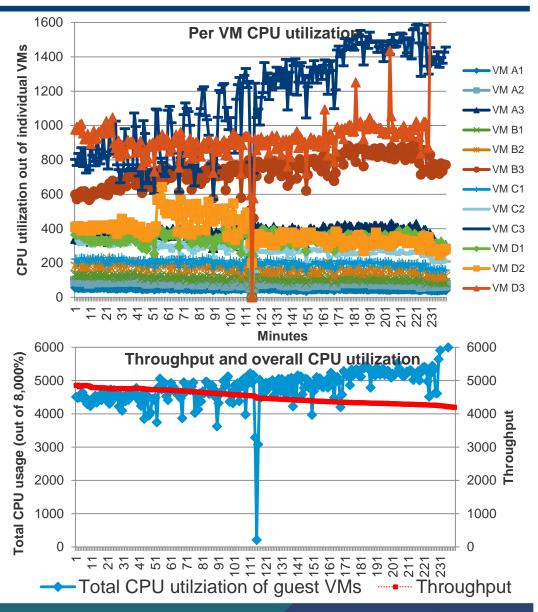
#### Carve out the 80-thread, 1TB server into:

- 4 Groups, each with 1 Set of 3 VMs
- 12 VMs
  - 4 Tier A VMs (VM1s of Groups A-D) with 2GB of memory and 3-8 vCPUs
    - These Tier A VMs have low resource demands
  - 4 DSS Tier B VMs (VM2s of Groups A-D) with 88-278GB of memory and 4-16 vCPUs
    - High I/O load (hence more memory to cache more of the database)
  - 4 OLTP Tier B VMs (VM3s of Groups A-D) with GB of 39-78GB of memory and 12-40 vCPUs
    - Low I/O load but high CPU demands
- Overall CPU allocation is overcommitted by 2X (168 vCPUs)
- No memory overcommit
  - Not recommended for database VMs

### Throughput and CPU utilization of 12 VMs with static load

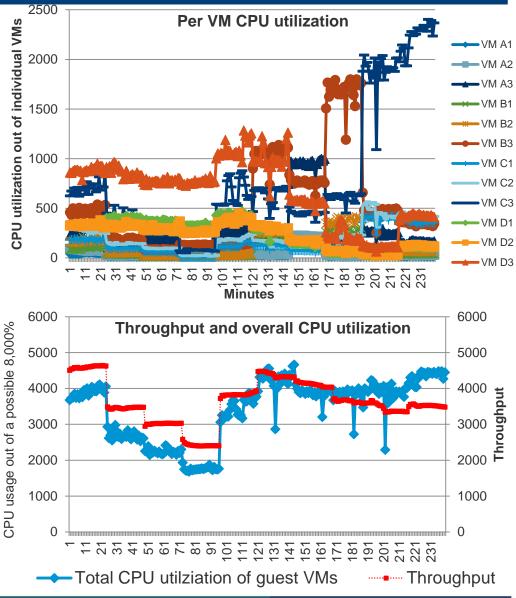
#### No elasticity

- Static load to all VMs over the 2 hours
- Throughput droped with time as CPU util% rises
  - Something wrong with VM C3. It utilization keeps rising, even above VM D3, despite static load
  - Early results. Our databases are better tuned now
- Server 85-95% utilized



### Throughput and CPU utilization of 12 VMs with elasticity

- Hypervisor has to react to 16X variation in load
  - Change how much CPU is allocated to each VM
- Overall CPU util% and throughput matches the static case in Phases 1 and 6 (no I/O bottleneck)
- The dips are the sign of the benchmark doing its job!!
  - Caused by storage for a VM getting overwhelmed
  - When a VM's storage can't keep up, its throughput drops
  - Kit maintains the ratios between all Groups, so overall throughput drops



**VLDB 2013** 

### Status of the kit and of the benchmark

- Functional specification in good shape, waiting for kit completion
- End-to-end kit running complete TPC-E workload
- Kit running all novel TPC-V functions, waiting for rewriting 2 transactions
- Would like to see more internal TPC prototyping before releasing to the wild
- Much PostgreSQL tuning remains
  - Both CPU cycles and IOPS/tran are over 2X of commercial databases
  - Will engage the community
  - Is the community interested in matching commercial databases?
- Will consider engaging groups outside TPC in the development process
  - Not typical for the TPC
  - But we will make the case for it if we can demonstrate serious commitment

### **Future directions**

### Complete the kit!

• Make it available for prototyping

#### Complete the spec

- Kit will become publicly available when the benchmark is released
- Need to release this benchmark in 2014
- Single-system virtualization no longer exciting!

#### Can we model a cloud-like, multi-server config?

- The benchmark and the kit already deal with elastic load to 1 server. Extend this to elasticity and migrations across multiple servers
- Intuitively, a simple step from where we are
- But need to come up with migration scenarios w/o too much choreographing
- Avoid deep pocket escalation wars with bigger and bigger clusters
- Add deployment and provisioning

# **Questions?**



**VLDB 2013**